

QUADRATIC ISOMETRIC FUNCTIONS OF THE HEISENBERG GROUPS. A COMBINATORIAL PROOF

A. Yu. Olshanskii and M. V. Sapir

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Quadratic isoperimetric inequalities for Heisenberg groups \mathcal{H}^{2n+1} have been announced in [2] and [4] and proved in [1]. The proof in [1] uses techniques and facts from topology and differential geometry. Here we present a completely combinatorial proof. The problem of finding such a proof was mentioned in [1].

Notice that \mathcal{H}^{2n+1} is not an automatic group [2]. Therefore it is not hyperbolic, and so it does not have a subquadratic isoperimetric function ([3], [5]).

Theorem 1. *The 5-dimensional Heisenberg group $\mathcal{H}^5 = \langle x, y, u, v | [x, u] = [x, v] = [y, u] = [y, v] = [x, y][v, u] = 1 \rangle$ has a quadratic isoperimetric function.*

This theorem and its proof can be easily extended to any Heisenberg group \mathcal{H}^{2n+1} with $n \geq 2$ and some other 2-nilpotent groups which are central products.

The following lemma contains some well-known and obvious facts about \mathcal{H}^5 . These facts will be used later without reference.

Lemma 2. 1. \mathcal{H}^5 is nilpotent of class 2.

2. The derived subgroup $(\mathcal{H}^5)'$ is infinite cyclic.

3. The factor group $\mathcal{H}^5/(\mathcal{H}^5)'$ is free Abelian of rank 4.

4. For every two elements $a, b \in \mathcal{H}^5$ and any two integers m, n , we have $[a^m, b^n] = [a, b]^{mn}$.

Definition 3. A word w is called a commutator word if it belongs to the derived subgroup of the free group. The notation $w(a, b)$ will mean that the word w depends on a, b only. If we substitute $a \rightarrow p$, $b \rightarrow q$ (where p, q are words), then the result of the substitution will be denoted by $w(p, q)$.

Lemma 4. In \mathcal{H}^5 , we have $(xu)^{v^{\pm 1}} = (xu)^{v^{\pm 1}}$, that is, $[xu, y^{\pm 1}v^{\mp 1}] = 1$.

(We assume that $a^b = b^{-1}ab$, $[a, b] = a^{-1}b^{-1}ab$.)

Proof. $(xu)^y = x^y u = x[x, y]u = x[u, v]u = xu[u, v] = xu^v = (xu)^v$. This implies that $(xu)^{y^{-1}} = (xu)^{v^{-1}}$ since $[y, v] = 1$. \square

The next lemma immediately follows from Lemma 4.

Lemma 5. Every word of the form $y^{\pm 1}w(xu, v)y^{\mp 1}$ can be transformed into the word $v^{\pm 1}w(xu, v)v^{\mp 1}$ by using $O(|w|)$ applications of relations.

Proof. $w(xu, v)^{y^{\pm 1}} = w((xu)^{y^{\pm 1}}, v^{y^{\pm 1}}) = w((xu)^{v^{\pm 1}}, v^{v^{\pm 1}}) = w(xu, v)^{v^{\pm 1}}$. \square

Lemma 6. Every commutator word $w(x, y)$ can be transformed into the word $w(u, v)$ by using at most $C_0|w|^2$ applications of defining relations of \mathcal{H}^5 for some $C_0 > 0$. In particular, $w(x, y) = w(u, v)$ in \mathcal{H}^5 .

Proof. Let n be the number of occurrences of the letter x in $w(x, y)$. Then the number of occurrences of x^{-1} is also n since w is a commutator word. Then we have

$$w(x, y) = w(x, y)u^n u^{-n} = w(xu, y).$$

Translated from Itogi Nauki i Tekhniki, Seriya Sovremennaya Matematika i Ee Prilozheniya. Tematicheskie Obzory. Vol. 55, Algebra-11, 1998.

The last transformation requires $O(n|w|) = O(|w|^2)$ applications of commutation relations $[x, u] = [y, u] = 1$.

Our goal is to replace the letters $y^{\pm 1}$ in $w(xu, y)$ by $v^{\pm 1}$ by replacing subwords of the form $y^{\pm 1}p(xu, v)y^{\mp 1}$ by $v^{\pm 1}p(xu, v)v^{\mp 1}$ using Lemma 5 and the fact that the numbers of occurrences of y and y^{-1} in $w(x, y)$ coincide. Note that although $w(xu, y)$ does not contain v , this letter will appear in it after the first application of Lemma 5.

The number of such transformations is the number of occurrences of y in $w(x, y)$. Each transformation requires $O(|w|)$ relations, so the total number of relations needed is $O(|w|^2)$. \square

Lemma 7. *Let $a, b \in \{x, x^{-1}, y, y^{-1}\}$. Then any word $w(a, b)$ of the form*

$$w_0 w_1 \dots w_s \prod_{i=1}^s [a, b]^{q_i}$$

where w_i, q_i are words and $|q_i| \leq Q$ for some number Q , can be transformed into the word

$$w_0 [x, y] w_1 [x, y] \dots w_s$$

by using $O((|w| + Q)s)$ applications of relations.

Proof. First we transform in $\prod_{i=1}^s [a, b]^{q_i}$ each $[a, b]$ into $[u^{\pm 1}, v^{\pm 1}]$, then move the commutator number $i = 1, 2, \dots$ to its spot between w_{i-1} and w_i (this requires $O(Q + |w|)$ commutativity relations for each i since q_1, \dots, q_{i-1} disappear when we start moving the i th commutator), and then transform each commutator $[u^{\pm 1}, v^{\pm 1}]$ back into $[a, b]$. \square

Definition 8. For every integer m, n, ℓ , $0 \leq \ell < m$, $m, n > 0$, let $R(m, n, \ell)$ be the following word $[x^m, y^n][x^\ell, y]^{y^n}$. If we draw this word on the plane R^2 (x -edges are horizontal, y -edges are vertical), then the image will be a rectangle with sides m and n with a horizontal ℓ by 1 rectangle on top. We define the *area* of the word $R(m, n, \ell)$ as $mn + \ell$. Note that by Lemma 2 $R(m, n, \ell) = [x, y]^{mn+\ell}$ in \mathcal{H}^5 and that the length of $R(m, n, \ell)$ in the free group is $2(m + n) + 2$.

The following statement is obvious.

Lemma 9. *For every integer m, n, ℓ , $0 < \ell \leq m$, the word $R(m, n, 0) \prod_{i=1}^\ell [x, y]^{q_i}$, where the length of the words q_i does not exceed some number Q , can be transformed into $R(m, n, \ell)$ if $\ell < m$ or $R(m, n + 1, 0)$ if $\ell = m$. The number of applications of relations used in this transformation is $C_1(m + n + Q)\ell$ for some constant $C_1 > 0$.*

Proof. Indeed, it is enough to divide $R(m, n, 0)$ into $\ell + 1$ subwords w_0, \dots, w_ℓ as

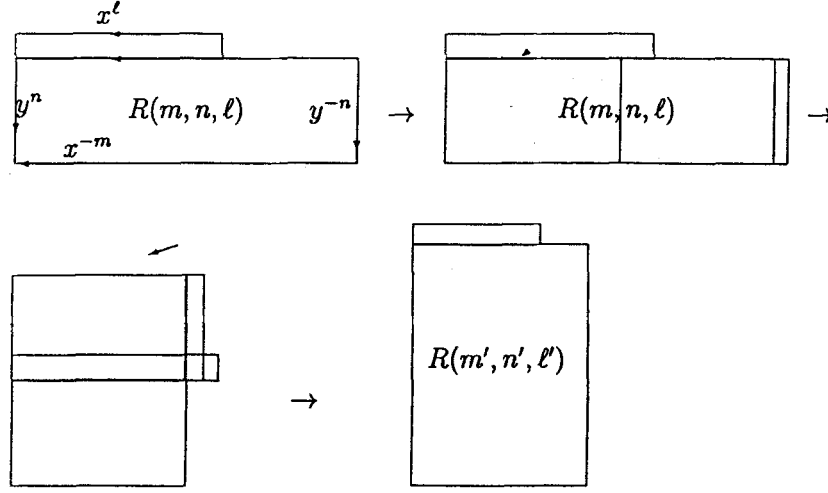
$$(x^{-m}y^{-n}x^{m-\ell+1}), x, x, \dots, x, y^n$$

and use Lemma 7. The resulting word will be equal to $R(m, n, \ell)$ (if $\ell < m$) or $R(m, n + 1, 0)$ ($\ell = m$) in the free group. \square

Lemma 10. *Let $a, b \in \{x, x^{-1}, y, y^{-1}\}$. Then in the free group, $[a^k, b] = \prod_{i=1}^k [a, b]^{a^{k-i}}$, $[a, b^k] = \prod_{i=0}^{k-1} [a, b]^{b^i}$.*

Lemma 11. *If $m > 3n$, then the word $R(m, n, \ell)$ can be transformed into a word $R(m', n', \ell')$ of the same area with $m' \leq m/2$ and $n' \leq 3m'$ by using $\leq C_2 m^2$ defining relations for some constant $C_2 > 0$. In addition, $m' + n' \leq m + n$.*

Proof



The idea of the proof is simple: we cut the rectangle approximately into two halves by a vertical line, then using Lemma 6 replace letters x and y in the right part by u and v , then move the right part on top of the left part using the commutativity relations, then replace u and v back into x and y , and finally distribute the small amount of extra squares to obtain a word of the form $R(m', n', \ell')$.

Let $m' = \lfloor m/2 \rfloor$, $t = m - 2m'$. Note that $t = 0$ or 1 . We have that

$$\begin{aligned}
 R(m, n, \ell) &= (\text{definition}) \\
 [x^m, y^n][x^\ell, y]^{y^n} &= (\text{in the free group}) \\
 [x^{m-m'}, y^n]^{x^{m'}}[x^{m'}, y^n][x^\ell, y]^{y^n} &= (\text{by Lemma 6, } O((m+n)^2) = O(m^2) \text{ relations needed}) \\
 [u^{m-m'}, v^n]^{x^{m'}}[x^{m'}, y^n][x^\ell, y]^{y^n} &= (O(m^2) \text{ commutativity relations}) \\
 [x^{m'}, y^n][u^{m-m'}, v^n][x^\ell, y]^{y^n} &= (O(m^2) \text{ commutativity relations}) \\
 [x^{m'}, y^n][u^{m-m'}, v^n]^{y^n}[x^\ell, y]^{y^n} &= (\text{by Lemma 6, } O(m^2) \text{ relations}) \\
 [x^{m'}, y^n][x^{m-m'}, y^n]^{y^n}[x^\ell, y]^{y^n} &= (\text{in the free group}) \\
 x^{-m'}y^{-n}x^{2m'-m}y^{-n}x^{m-m'}y^{2n}[x^\ell, y]^{y^n} &= (\text{in the free group}) \\
 [x^{m'}, y^{2n}][y^{-n}, x^t]^{x^{m'}}y^{2n}[x^\ell, y]^{y^n} &
 \end{aligned}$$

By Lemma 10, $[y^{-n}, x^t]$ is equal to $\prod_{i=1}^n [y^{-1}, x^t]^{y^{-(n-i)}}$ and $[y^{-1}, x^t] = [x^t, y]^{y^{-1}}$ in the free group. Again by Lemma 10, $[x^\ell, y] = \prod_{i=1}^\ell [x, y]^{x^{\ell-i}}$. Therefore

$$[x^{m'}, y^{2n}][y^{-n}, x^t]^{x^{m'}}y^{2n}[x^\ell, y]^{y^n} = [x^{m'}, y^{2n}] \prod_{i=1}^s [x, y]^{q_i} = R(m', 2n, 0) \prod_{i=1}^s [x, y]^{q_i}$$

where $|q_i| = O(m)$, $s \leq n + \ell$. Note that

$$s \leq n + \ell < m/3 + 2m' \leq (2m' + 1)/3 + 2m' < 3m'.$$

Since $s \leq 3m'$, using Lemma 9 at most three times, one can transform the word $R(m', 2n, 0) \prod_{i=1}^s [x, y]^{q_i}$ into a word $R(m', n', \ell')$ for some n', ℓ' where $m'n' + \ell' = mn + \ell$ and $n' \leq 2n + 2$.

It is clear that $m' \leq m/2$ and $n' \leq 3m'$ since $n' < m/3 + 2$. Finally $m' + n' \leq m/2 + 2n + 2 \leq m + n$: this is clearly true if $m \geq 10$ (since $m > 3n$) and can be easily checked for every $m < 10$ (for example, if $m = 7, n = 2, \ell = 6$, then $m' = 3, n' = 6, \ell' = 2$, and $m' + n' = m + n = 9$). \square

The following lemma has exactly the same proof.

Lemma 12. *If $n > 3m$, then the word $R(m, n, \ell)$ can be transformed into a word $R(m', n', \ell')$ of the same area with $n' \leq n/2$ and $m' \leq 3n'$ by using $\leq C_2 n^2$ defining relations for some constant $C_2 > 0$. In addition, $m' + n' \leq m + n$.*

We call the word $R(m, n, \ell)$ *thick* if both $m \leq 3n$ and $n \leq 3m$.

Lemma 13. *Every word $R(m, n, \ell)$ can be transformed into a thick word $R(m_0, n_0, \ell_0)$ of the same area by using at most $C_3(m + n)^2$ defining relations for $C_3 = \frac{4}{3}C_2$. In addition, $m_0 + n_0 \leq m + n$.*

Proof. Without loss of generality, assume that $m > 3n$. Then by Lemma 11 we can transform the word $R(m, n, \ell)$ into the word $R(m', n', \ell')$ of the same area with $m' \leq m/2$, $n' \leq 3m'$ by using at most $C_2 m^2$ defining relations. If $m' > 3n'$, we can apply Lemma 11 again and get a word $R(m'', n'', \ell'')$ with $m'' \leq m'/2 \leq m/4$, $n'' \leq 3m''$ by using at most $C_2(m')^2 \leq C_2 m^2/4$. Continuing in this manner, we finally get the desired thick word $R(m_0, n_0, \ell_0)$ by applying less than $C_2 m^2(1 + 1/4 + 1/16 + \dots) \leq C_3 m^2$ defining relations. The last statement of the lemma follows from Lemmas 11 and 12. \square

Lemma 14. *Every two thick words $R_1 = R(m_1, n_1, \ell_1)$ and $R_2 = R(m_2, n_2, \ell_2)$ of the same area A can be transformed into one another by using at most $C_4 A$ defining relations for some constant $C_4 > 0$.*

Proof. We use induction on the area A . Note that if $m_1 = m_2$, then the words R_1 and R_2 are graphically equal since they have the same area $A = m_1 n_1 + \ell_1 = m_2 n_2 + \ell_2$, $\ell_1 < m_1$, $\ell_2 < m_2$.

So without loss of generality, assume that $m_1 > m_2$. Then $n_1 \leq n_2$. Since both words R_1 and R_2 are thick, it is easy to check that each of the numbers m_1, n_1, m_2, n_2 is at least $1/3$ of any other number in this list. This implies that each of these numbers is at most $2\sqrt{A}$.

We have that in the free group

$$R_1 = R(m_1 - m_2, n_1, 0)^{x^{m_2}} R(m_2, n_1, 0)[x^{\ell_1}, y]^{y^{n_1}}.$$

By Lemma 6 we can change the letters in $R(m_2, n_1, 0)$ from x, y to u, v , commute the resulting word with $R(m_1 - m_2, n_1, 0)^{x^{m_2}}$, and then convert the letters u, v back to x and y by applying at most $8C_0(m_2 + n_1)^2 + 4(m_2 + n_1)(m_1 + n_1) \leq D_1 A$ (for some constant D_1) defining relations. We get the word

$$R(m_2, n_1, 0)R(m_1 - m_2, n_1, 0)^{x^{m_2}}[x^{\ell_1}, y]^{y^{n_1}}.$$

Now applying Lemma 6 to $R(m_1 - m_2, n_1, 0)$, commuting with x^{m_2} , then replacing u, v with x, y (Lemma 6 again), we can remove the conjugation by x^{m_2} by using at most $D_2 A$ defining relations for some constant $D_2 > 0$. The word we get will have the form

$$R(m_2, n_1, 0)R(m_1 - m_2, n_1, 0)[x^{\ell_1}, y]^{y^{n_1}}.$$

Let $\ell_1 = (m_1 - m_2)r + \ell'$, where $0 \leq \ell' < m_1 - m_2$. Applying Lemma 10 to $[x^{\ell_1}, y]$ and then applying Lemma 9 r or $r + 1$ times (depending on whether $\ell' = 0$ or not), one can transform the word $R(m_1 - m_2, n_1, 0)[x^{\ell_1}, y]^{y^{n_1}}$ into $R(m_1 - m_2, n', \ell')$, where $n' = n_1 + r < n_1 + m_1$, by using at most

$$C_1(m_1 - m_2 + n_1 + m_1 + n_1)((m_1 - m_2)r + \ell') = C_1(m_1 - m_2 + n_1 + m_1 + n_1)\ell_1 \leq D_3 A$$

applications of defining relations. The resulting word has the form

$$R(m_2, n_1, 0)R(m_1 - m_2, n', \ell').$$

Now by Lemma 13 we can transform the second factor of this product into a thick word $R(m_3, n_3, \ell_3)$ with the same area and $m_3 + n_3 \leq m_1 - m_2 + n'$ by using $C_3(m_1 - m_2 + n')^2 \leq D_4 A$ defining relations for some constant D_4 .

The area of $R(m_3, n_3, \ell_3)$ is $A - m_2 n_1 = m_1 n_1 + \ell_1 - m_2 n_1 < \frac{5}{6}A$ since $m_2 \geq \frac{1}{3}m_1$ and $\ell_1 < m_1 n_1$.

Now let us deal with $R(m_2, n_2, \ell_2)$. We can rewrite it in the free group as follows:

$$R_2 = R(m_2, n_1, 0)R(m_2, n_2 - n_1, \ell_2)^{y^{n_1}}$$

if $n_2 > n_1$, or $R_2 = R(m_2, n_1, 0)$ if $n_1 = n_2$.

By using Lemma 6 we can remove (as before) the conjugation by y^{n_1} using at most $D_2 A$ defining relations. After that we can transform the word $R(m_2, n_2 - n_1, \ell_2)$ into a thick word $R(m_4, n_4, \ell_4)$ of the same area by using at most $D_4 A$ defining relations.

The area of $R(m_4, n_4, \ell_4)$ is equal to $A - m_2 n_1$, so it is equal to the area of $R(m_3, n_3, \ell_3)$ and is at most $\frac{5}{6}A$.

Now let us denote the constant $6(D_1 + 2D_2 + D_3 + 2D_4)$ by C_4 and assume by induction that every two thick words of area $B < A$ can be transformed into one another by applying at most $C_4 B$ defining relations.

Then we can transform the word $R(m_3, n_3, \ell_3)$ into $R(m_4, n_4, \ell_4)$ by using at most $\frac{5}{6}C_4 A$ defining relations. This implies that we can transform R_1 into R_2 by applying at most $D_1 A + D_2 A + D_3 A + D_4 A + D_2 A + D_4 A + \frac{5}{6}C_4 A = C_4 A$ defining relations. \square

Lemma 15. *For any two words $R(m_1, n_1, \ell_1)$ and $R(m_2, n_2, \ell_2)$ of the same area, one can transform one word into another by using at most $C_5(m_1 + n_1 + m_2 + n_2)^2$ defining relations for some constant $C_5 > 0$.*

Proof. This follows immediately from Lemmas 13 and 14: first we transform both words into thick words of the same area, then transform one thick word into another; we use the fact that when we transform a word $R(m, n, \ell)$ into a thick word $R(m', n', \ell')$ then $m' + n' \leq m + n$ (Lemma 13) and the trivial observation that the area of a word $R(m, n, \ell)$ does not exceed $(m + n)^2$. \square

Lemma 16. *For every two words $R_1 = R(m_1, n_1, \ell_1)$ and $R_2 = R(m_2, n_2, \ell_2)$ of areas A_1 and A_2 , their product $R_1 R_2$ can be transformed into a thick word $R(m_3, n_3, \ell_3)$ with area $A_1 + A_2$ by using at most $C_6(m_1 + n_1 + m_2 + n_2)^2$ defining relations for some constant $C_6 > 0$.*

Proof. We shall consider only the case where $A_1 \geq A_2$. The other case is similar. By Lemma 15 we can transform R_2 into a word R' where either $R' = R(m_1, n', \ell')$ (if $A_2 \geq m_1$) or $R' = R(A_2, 1, 0)$ (if $A_2 < m_1$) by using at most $O((m_1 + n_1 + m_2 + n_2)^2)$ defining relations.

As before, we consider only the most difficult case where $A_2 \geq m_1$ (in the other case the number of transformations will be smaller). We have

$$R_1 R' = R(m_1, n_1, 0)[x^{\ell_1}, y]^{y^{n_1}} R(m_1, n', 0)[x^{\ell'}, y]^{y^{n'}}$$

Using Lemma 6 we can permute $[x^{\ell_1}, y]^{y^{n_1}}$ and $R(m_1, n', 0)$, and replace $R(m_1, n', 0)$ by $R(m_1, n', 0)^{y^{n_1}}$ (we lift the second rectangle onto the top of the first rectangle). Note that $R(m_1, n_1, 0)R(m_1, n', 0)^{y^{n_1}} = R(m_1, n_1 + n', 0)$ in the free group. Then, using Lemmas 9 and 10, we convert the word

$$R(m_1, n_1 + n', 0)[x^{\ell_1}, y]^{y^{n_1}}[x^{\ell'}, y]^{y^{n'}}$$

into the word $R(m_1, n_3, \ell_3)$ for some n_3 and ℓ_3 . Now we can make this word thick by using Lemma 13. All these transformations require application of $O((m_1 + n_1 + m_2 + n_2)^2)$ defining relations. \square

Lemma 17. *For every two words $R_1 = R(m_1, n_1, \ell_1)$ and $R_2 = R(m_2, n_2, \ell_2)$ of areas A_1 and A_2 , their product $R_1 R_2^{-1}$ can be transformed into an empty word or a word $R^{\pm 1} = R(m_3, n_3, \ell_3)^{\pm 1}$ such that R is thick and the area of R is $|A_1 - A_2|$ by using at most $C_7(m_1 + n_1 + m_2 + n_2)^2$ defining relations for some constant $C_7 \geq C_6$.*

Proof. If $A_1 = A_2$, then we apply Lemma 15. So let $A_1 \neq A_2$. Without loss of generality assume that $A_1 > A_2$. Consider a thick word $R_3 = R(m_3, n_3, \ell_3)$ of area $A_1 - A_2$. Then $m_3 + n_3 = O(m_1 + n_1 + m_2 + n_2)$ (since R_3 is thick). By Lemma 16, we can transform $R_3 R_2$ into a thick word $R' = R(m', n', \ell')$ with area A_1 and $m' + n' = O(m_1 + n_1 + m_2 + n_2)$. By Lemma 15, we can transform R' into R_1 . The total number of applications of defining relations in all these transformations is $O((m_1 + n_1 + m_2 + n_2)^2)$. Thus, one can transform $R_1 R_2^{-1}$ into $R_3 R_2 R_2^{-1} = R_3$ by using $O((m_1 + n_1 + m_2 + n_2)^2)$ defining relations. \square

Lemma 18. *Every commutator word of the form $W = x^a w(x, y) y^b$ can be transformed either into the empty word or into a word $R(m, n, \ell)^{\pm 1}$ for some m, n, ℓ such that $R(m, n, \ell)$ is a thick word and its area, $mn + \ell$, is at most $|w|^2$ by using $C_8 |w|^2$ defining relations, where $C_8 = 250C_0 + 70C_7$.*

Proof. Induction on the length of w . The base ($|w| = 1$) is obvious.

Let us write w as $w_1 w_2$, where $|w_1| = \lfloor |w|/2 \rfloor$. Let p (resp. q) be the sum of exponents of x (resp. y) in w_1 . Then in the free group, we have

$$W = (x^{a+p} x^{-p} w_1 y^{-q} x^{-(a+p)}) (x^{a+p} y^q x^{-(a+p)} y^{-q}) (y^q x^{a+p} w_2 y^{b+q} y^{-q}).$$

Note that $|a + p|$ does not exceed the number of occurrences of x in w_2 and $|q|$ does not exceed the number of occurrences of y in w_1 . Therefore $|a + p| + |q| \leq |w|$.

Note that $x^{-p} w_1 y^{-q}$ and $x^{a+p} w_2 y^{b+q}$ are commutator words by the definition of p and q . By the induction hypothesis these words can be converted to the empty word or words of the form $R_i^{\pm 1} = R(s_i, t_i, \ell_i)^{\pm 1}$ for some s_i, t_i, ℓ_i , $i = 1, 2$, where R_i is a thick word of area at most $|w_i|^2$. The number of defining relations needed for these transformations is $C_8(|w_1|^2 + |w_2|^2) \leq \frac{5}{9} C_8 |w|^2$ (the worst case here is when $|w| = 3$, $|w_1| = 1$, $|w_2| = 2$).

Since R_i , $i = 1, 2$, is a thick word, $s_i + t_i \leq 3|w_i|$ (indeed, $3s_i \geq t_i$, $3t_i \geq s_i$, and $s_i t_i \leq |w_i|^2$).

Therefore, by using Lemma 6 we can remove conjugations by x^{a+p} and y^{-q} . The number of defining relations needed for these transformation is at most $100C_0 |w|^2$ (the length of R_i is at most $2 \cdot 3|w_i| + 2 \leq 8|w_i|$; we need to replace x, y by u, v , then commute resulting words with the conjugating elements, then convert x, y back to u, v).

The commutator $x^{a+p} y^q x^{-(a+p)} y^{-q}$ is conjugated in the free group to a word of the form $R_3^{\pm 1} R(|a + p|, |q|, 0)^{\pm 1}$, and the conjugating word has length at most $|w|$. This conjugating word can be removed by using Lemma 6, which takes at most $10C_0 |w|^2$ defining relations. The area of R_3 is $|a + p||q| \leq |w|^2/4$ since $|a + p| + |q| \leq |w|$.

The resulting word is a product of three words $R_1^{\pm 1} R_3^{\pm 1} R_2^{\pm 1}$. Applying Lemmas 16 and 17, we can combine these words into a word $R(m, n, \ell)^{\pm 1}$. This operation takes at most $30C_7 |w|^2$.

The area of $R(m, n, \ell)$ is at most the sum of the areas of R_1 , R_2 , and R_3 (by Lemmas 16 and 17). Therefore, the area of $R(m, n, \ell)$ is at most $\frac{5}{9} |w|^2 + \frac{1}{4} |w|^2 \leq |w|^2$ (as desired).

The total number of defining relations needed in order to get $R(m, n, \ell)^{\pm 1}$ is at most

$$(\frac{5}{9} C_8 + 110C_0 + 30C_7) |w|^2 \leq C_8 |w|^2. \quad \square$$

Proof of the Theorem. Let $w = w(x, y, u, v)$ be any word in x, y, u, v which is equal to 1 in \mathcal{H}^5 . Using $O(|w|^2)$ commutativity relations, we can transform w into a word of the form $w_1(x, y) w_2(u, v)$ which has the same length as w . By part 3 of Lemma 2, both w_1 and w_2 are commutator words. Therefore, by Lemma 6 we can transform $w_2(u, v)$ into $w_2(x, y)$ by using $O(|w|^2)$ defining relations. Thus we can assume that $w = w(x, y)$.

By Lemma 18, for $a = b = 0$ we can transform the word w into a word $R^{\pm 1} = R(m, n, \ell)^{\pm 1}$ or the empty word by using $O(|w|^2)$ defining relations. By Lemma 2, the word R is equal to $[x, y]^{mn+\ell} \neq 1$ in \mathcal{H}^5 because $[x, y]$ is an element of infinite order and $m, n > 0, \ell \geq 0$. Therefore, our transformation changes w into the

empty word. Therefore, w can be transformed into the empty word by using $O(|w|^2)$ defining relations. Thus, \mathcal{H}^5 has a quadratic isoperimetric function. \square

Remark. The method used in this paper can be employed to find isoperimetric functions of other nilpotent groups. For example, one can prove that the central product of $l \geq 2$ copies of any finitely generated free nilpotent group of class 2 with derived subgroups identified in the natural way has a quadratic isoperimetric function. We can also prove that such a central product of two copies of any finitely generated nilpotent group of class 2 has an isoperimetric function equivalent to $n^2 \log n$. So far, we have been unable to prove that these groups always have quadratic isoperimetric functions.

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